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## Efficient Algorithms and Datastructures II

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### Aufgabe 1 (10 Punkte)

Let  $G = (V, E)$  be a given graph and  $c_e \geq 0$  be the cost of edge  $e$ . Let  $\{(s_1, t_1), \dots, (s_k, t_k)\}$  be a set of specified pairs of vertices. In the minimum multicut problem, we wish to find a minimum cost set of edges  $F$  such that  $\forall i, s_i$  and  $t_i$  are in different components of  $G' = (V, E \setminus F)$ .

- Write an Integer Linear Program (ILP) for solving this problem, where you have a variable for each edge and a constraint for each path from  $s_i$  to  $t_i$ , for all  $i$ .
- Relax this ILP to a Linear Program, say (P).
- Show how to solve (P) efficiently.

### Aufgabe 2 (10 Punkte)

Given a directed graph  $G = (V, E)$ , a special vertex  $r$  and a positive cost  $c_{ij}$  for each edge  $(i, j) \in E$ , the minimum-cost arborescence problem is to find a subgraph of minimum cost that contains directed paths from  $r$  to all other vertices.

- Observe that the following ILP solves the minimum-cost arborescence problem:

$$\begin{aligned} & \text{minimize} && \sum_{(i,j) \in E} c_{ij} x_{ij} \\ & \text{subject to} && \sum_{i \in S, j \notin S, (i,j) \in E} x_{ij} \geq 1 \quad \forall S \subseteq V, S \ni r \\ & && x_{ij} \in \{0, 1\} \quad \forall (i, j) \in E \end{aligned}$$

- Show how to efficiently solve the LP obtained by relaxing the above ILP.

### Aufgabe 3 (10 Punkte)

Let  $G = (V, E)$  be a given graph. Consider the following ILP:

$$\begin{aligned} & \text{maximize} && \sum_i x_i \\ & \text{subject to} && x_i + x_j \leq 1 \quad \forall (i, j) \in E \\ & && \sum_{i \in C} x_i \leq \frac{|C|-1}{2} \quad \forall \text{ odd cycles } C \\ & && x_i \in \{0, 1\} \quad \forall i \in V \end{aligned}$$

- Explain in your own words, which problem the above ILP solves.
- Relax this ILP to an LP so that  $0 \leq x_i \leq 1, \forall i \in V$  and show how to solve this LP efficiently.