- inequalities $Ax \leq b$; $m \times n$ matrix A with rows a_i^T
- $P = \{x \mid Ax \le b\}; P^{\circ} := \{x \mid Ax < b\}$
- interior point algorithm: $x \in P^\circ$ throughout the algorithm
- for $x \in P^\circ$ define

$$s_i(x) := b_i - a_i^T x$$

as the slack of the *i*-th constraint

logarithmic barrier function:

$$\phi(x) = -\sum_{i=1}^m \log(s_i(x))$$

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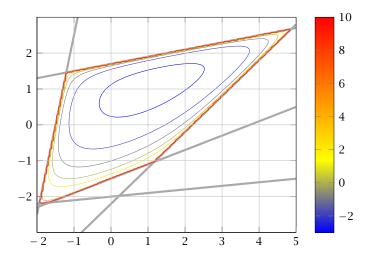
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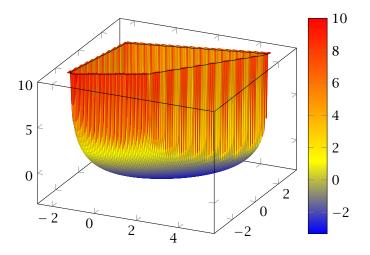
Penalty Function





10 Karmarkars Algorithm

Penalty Function





Gradient and Hessian

Taylor approximation:

$$\phi(x+\epsilon) \approx \phi(x) + \nabla \phi(x)^T \epsilon + \frac{1}{2} \epsilon^T \nabla^2 \phi(x) \epsilon$$

Gradient:

$$\nabla \phi(x) = \sum_{i=1}^{m} \frac{1}{s_i(x)} \cdot a_i = A^T d_x$$

where $d_x^T = (1/s_1(x), \dots, 1/s_m(x))$. (d_x vector of inverse slacks)

Hessian:

$$H_{\mathbf{x}} := \nabla^2 \phi(\mathbf{x}) = \sum_{i=1}^m \frac{1}{s_i(\mathbf{x})^2} a_i a_i^T = A^T D_{\mathbf{x}}^2 A_i$$

with $D_x = \operatorname{diag}(d_x)$.

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with $D_x = \text{diag}(d_x)$.

Proof for Gradient

$$\begin{split} \frac{\partial \phi(x)}{\partial x_i} &= \frac{\partial}{\partial x_i} \left(-\sum_r \ln(s_r(x)) \right) \\ &= -\sum_r \frac{\partial}{\partial x_i} \left(\ln(s_r(x)) \right) = -\sum_r \frac{1}{s_r(x)} \frac{\partial}{\partial x_i} \left(s_r(x) \right) \\ &= -\sum_r \frac{1}{s_r(x)} \frac{\partial}{\partial x_i} \left(b_r - a_r^T x \right) = \sum_r \frac{1}{s_r(x)} \frac{\partial}{\partial x_i} \left(a_r^T x \right) \\ &= \sum_r \frac{1}{s_r(x)} A_{ri} \end{split}$$

The *i*-th entry of the gradient vector is $\sum_{r} 1/s_r(x) \cdot A_{ri}$. This gives that the gradient is

$$\nabla \phi(x) = \sum_{r} 1/s_{r}(x)a_{r} = A^{T}d_{x}$$

Proof for Hessian

$$\frac{\partial}{\partial x_j} \left(\sum_r \frac{1}{s_r(x)} A_{ri} \right) = \sum_r A_{ri} \left(-\frac{1}{s_r(x)^2} \right) \cdot \frac{\partial}{\partial x_j} \left(s_r(x) \right)$$
$$= \sum_r A_{ri} \frac{1}{s_r(x)^2} A_{rj}$$

Note that $\sum_{r} A_{ri}A_{rj} = (A^{T}A)_{ij}$. Adding the additional factors $1/s_{r}(x)^{2}$ can be done with a diagonal matrix.

Hence the Hessian is

$$H_X = A^T D^2 A$$

 H_X is positive semi-definite for $x \in P^\circ$

 $u^{T}H_{x}u = u^{T}A^{T}D_{x}^{2}Au = ||D_{x}Au||_{2}^{2} \ge 0$

This gives that $\phi(x)$ is convex.

If rank(A) = n, H_x is positive definite for $x \in P^\circ$ $u^T H_x u = \|D_x A u\|_2^2 > 0$ for $u \neq 0$

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Points in Ex are feasible!!!

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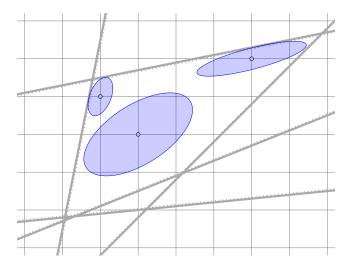
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$$\leq 1$$





10 Karmarkars Algorithm

Analytic Center

 $x_{\mathrm{ac}} := \operatorname{arg\,min}_{x \in P^\circ} \phi(x)$

• x_{ac} is solution to

$$\nabla \phi(x) = \sum_{i=1}^m \frac{1}{s_i(x)} a_i = 0$$

- depends on the description of the polytope
- x_{ac} exists and is unique iff P° is nonempty and bounded



In the following we assume that the LP and its dual are strictly feasible and that rank(A) = n.

```
Central Path:
Set of points \{x^*(t) \mid t > 0\} with
```

 $x^*(t) = \operatorname{argmin}_x \{ t c^T x + \phi(x) \}$

- t = 0: analytic center
- $t = \infty$: optimum solution

 $x^*(t)$ exists and is unique for all $t \ge 0$.



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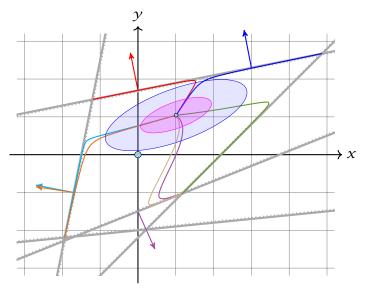
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Different Central Paths





10 Karmarkars Algorithm

Intuitive Idea:

Find point on central path for large value of t. Should be close to optimum solution.

Questions:

- Is this really true? How large a t do we need?
- How do we find corresponding point $x^*(t)$ on central path?



The Dual

primal-dual pair:

Assumptions

- primal and dual problems are strictly feasible;
- $\operatorname{rank}(A) = n$.

Force Field Interpretation

Point $x^*(t)$ on central path is solution to $tc + \nabla \phi(x) = 0$

- We can view each constraint as generating a repelling force. The combination of these forces is represented by ∇φ(x).
- In addition there is a force tc pulling us towards the optimum solution.



How large should *t* be?

Point $x^*(t)$ on central path is solution to $tc + \nabla \phi(x) = 0$.

This means

$$tc + \sum_{i=1}^{m} \frac{1}{s_i(x^*(t))} a_i = 0$$

or

$$c + \sum_{i=1}^{m} z_i^*(t) a_i = 0$$
 with $z_i^*(t) = \frac{1}{t s_i(x^*(t))}$

2000 is strictly dual feasible: (2022-5-0-0) 225-0)

if gap is less than $1/2^{1010}$ we can snap to optimum point:

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 duality gap between *x* := *x**(*t*) and *z* := *z**(*t*) is

$$c^T x + b^T z = (b - Ax)^T z = \frac{m}{t}$$

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How to find $x^*(t)$

First idea:

- start somewhere in the polytope
- use iterative method (Newtons method) to minimize $f_t(x) := tc^T x + \phi(x)$



Quadratic approximation of f_t

$$f_t(x + \epsilon) \approx f_t(x) + \nabla f_t(x)^T \epsilon + \frac{1}{2} \epsilon^T H_{f_t}(x) \epsilon$$

Suppose this were exact:

$$f_t(x + \epsilon) = f_t(x) + \nabla f_t(x)^T \epsilon + \frac{1}{2} \epsilon^T H_{f_t}(x) \epsilon$$

Then gradient is given by:

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10 Karmarkars Algorithm

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We want to move to a point where this gradient is 0:

Newton Step at $x \in P^{\circ}$

$$\Delta x_{\mathsf{nt}} = -H_{f_t}^{-1}(x)\nabla f_t(x)$$

= $-H_{f_t}^{-1}(x)(tc + \nabla \phi(x))$
= $-(A^T D_x^2 A)^{-1}(tc + A^T d_x)$

Newton Iteration:

$$x := x + \Delta x_{nt}$$

Measuring Progress of Newton Step

Newton decrement:

 $\lambda_t(x) = \|D_x A \Delta x_{\mathsf{nt}}\| \\ = \|\Delta x_{\mathsf{nt}}\|_{H_x}$

Square of Newton decrement is linear estimate of reduction if we do a Newton step:

 $-\lambda_t(x)^2 = \nabla f_t(x)^T \Delta x_{\mathsf{nt}}$

• $\lambda_t(x) = 0$ iff $x = x^*(t)$

• $\lambda_t(x)$ is measure of proximity of x to $x^*(t)$

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Theorem 2 If $\lambda_t(x) < 1$ then

- $x_+ := x + \Delta x_{nt} \in P^\circ$ (new point feasible)
- $\lambda_t(x_+) \leq \lambda_t(x)^2$

This means we have quadratic convergence. Very fast.

feasibility:

► $\lambda_t(x) = \|\Delta x_{nt}\|_{H_x} < 1$; hence x_+ lies in the Dikin ellipsoid around x.

bound on $\lambda_t(x^+)$: we use $D := D_x = \text{diag}(d_x)$ and $D_+ := D_{x^+} = \text{diag}(d_{x^+})$

To see the last equality we use Pythagoras

 $||a||^2 + ||a + b||^2 = ||b||^2$

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 $\lambda_{t}(\boldsymbol{x}^{+})^{2} = \|D_{+}A \Delta x_{nt}^{+}\|^{2}$ $\leq \|D_{+}A \Delta x_{nt}^{+}\|^{2} + \|D_{+}A \Delta x_{nt}^{+} + (I - D_{+}^{-1}D)DA \Delta x_{nt}\|^{2}$ $= \|(I - D_{+}^{-1}D)DA \Delta x_{nt}\|^{2}$

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If $\lambda_t(x)$ is large we do not have a guarantee.

Try to avoid this case!!!



Path-following Methods

Try to slowly travel along the central path.

Algorithm 1 PathFollowing

- 1: start at analytic center
- 2: while solution not good enough do
- 3: make step to improve objective function
- 4: recenter to return to central path

simplifying assumptions:

- a first central point $x^*(t_0)$ is given
- $x^*(t)$ is computed exactly in each iteration

ϵ is approximation we are aiming for

start at $t = t_0$, repeat until $m/t \le \epsilon$

- compute $x^*(\mu t)$ using Newton starting from $x^*(t)$
- ► *t* := *µt*

where $\mu = 1 + 1/(2\sqrt{m})$

gradient of f_{t^+} at ($x = x^*(t)$)

$$\nabla f_{t^+}(x) = \nabla f_t(x) + (\mu - 1)tc$$
$$= -(\mu - 1)A^T D_X \vec{1}$$

This holds because $0 = \nabla f_t(x) = tc + A^T D_x \vec{1}$.

The Newton decrement is

$$\begin{split} \lambda_{t^{+}}(x)^{2} &= \nabla f_{t^{+}}(x)^{T} H^{-1} \nabla f_{t^{+}}(x) \\ &= (\mu - 1)^{2} \vec{1}^{T} B (B^{T} B)^{-1} B^{T} \vec{1} \qquad B = D_{x}^{T} A \\ &\leq (\mu - 1)^{2} m \\ &= 1/4 \end{split}$$

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Number of Iterations

the number of Newton iterations per outer iteration is very small; in practise only 1 or 2

Number of outer iterations:

We need $t_k = \mu^k t_0 \ge m/\epsilon$. This holds when

 $k \geq \frac{\log(m/(\epsilon t_0))}{\log(\mu)}$

We get a bound of

$$\mathcal{O}\left(\sqrt{m}\log\frac{m}{\epsilon t_0}\right)$$

We show how to get a starting point with $t_0 = 1/2^L$. Together with $\epsilon \approx 2^{-L}$ we get $\mathcal{O}(L\sqrt{m})$ iterations.

EADS II Harald Räcke

For $x \in P^{\circ}$ and direction $v \neq 0$ define

$$\sigma_{x}(v) := \max_{i} \frac{a_{i}^{T} v}{s_{i}(x)}$$

Observation:

$$x + \alpha v \in P$$
 for $\alpha \in \{0, 1/\sigma_x(v)\}$



Suppose that we move from x to $x + \alpha v$. The linear estimate says that $f_t(x)$ should change by $\nabla f_t(x)^T \alpha v$.

The following argument shows that f_t is well behaved. For small α the reduction of $f_t(x)$ is close to linear estimate.

 $f_t(x+\alpha v)-f_t(x)=tc^T\alpha v+\phi(x+\alpha v)-\phi(x)$

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Define $w_i = a_i^T v / s_i(x)$ and $\sigma = \max_i w_i$. Then

 $f_t(x + \alpha v) - f_t(x) - \nabla f_t(x)^T \alpha v$

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Damped Newton Iteration: In a damped Newton step we choose

$$x_{+} = x + \frac{1}{1 + \sigma_x(\Delta x_{\rm nt})} \Delta x_{\rm nt}$$



10 Karmarkars Algorithm

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Theorem:

In a damped Newton step the cost decreases by at least

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Proof: The decrease in cost is

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 $\geq \lambda_t(x) - \log(1 + \lambda_t(x))$ ≥ 0.09

for $\lambda_t(x) \ge 0.5$

Centering Algorithm: Input: precision δ ; starting point x**1.** compute Δx_{nt} and $\lambda_t(x)$ **2.** if $\lambda_t(x) \le \delta$ return x**3.** set $x := x + \alpha \Delta x_{nt}$ with

$$\alpha = \begin{cases} \frac{1}{1 + \sigma_x(\Delta x_{\mathsf{nt}})} & \lambda_t \ge 1/2\\ 1 & \text{otw.} \end{cases}$$



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Centering

Lemma 3

The centering algorithm starting at x_0 reaches a point with $\lambda_t(x) \le \delta$ after

$$\frac{f_t(x_0) - \min_{\mathcal{Y}} f_t(\mathcal{Y})}{0.09} + \mathcal{O}(\log \log(1/\delta))$$

iterations.

This can be very, very slow...



Let $P = \{Ax \le b\}$ be our (feasible) polyhedron, and x_0 a feasible point.

We change $b \to b + \frac{1}{\lambda} \cdot \vec{1}$, where $L = \langle A \rangle + \langle b \rangle + \langle c \rangle$ (encoding length) and $\lambda = 2^{2L}$. Recall that a basis is feasible in the old LP iff it is feasible in the new LP.



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Lemma [without proof] The inverse of a matrix M can be represented with rational numbers that have denominators $z_{ij} = det(M)$.

For two basis solutions x_B , $x_{\bar{B}}$, the cost-difference $c^T x_B - c^T x_{\bar{B}}$ can be represented by a rational number that has denominator $z = \det(A_B) \cdot \det(A_{\bar{B}}) \cdot \lambda$.

This means that in the perturbed LP it is sufficient to decrease the duality gap to $1/2^{4L}$ (i.e., $t \approx 2^{4L}$). This means the previous analysis essentially also works for the perturbed LP.

For a point x from the polytope (not necessarily BFS) the objective value $\bar{c}^T x$ is at most $n2^M 2^L$, where $M \leq L$ is the encoding length of the largest entry in \bar{c} .



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Start at x_0 .

Choose $\hat{c} := -\nabla \phi(x)$.

 $x_0 = x^*(1)$ is point on central path for \hat{c} and t = 1.

You can travel the central path in both directions. Go towards 0 until $t \approx 1/2^{\Omega(L)}$. This requires $O(\sqrt{m}L)$ outer iterations.

Let $x_{\hat{c}}$ denote this point.

Let x_c denote the point that minimizes

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Clearly,

$$t \cdot \hat{c}^T \boldsymbol{x}_{\hat{c}} + \phi(\boldsymbol{x}_{\hat{c}}) \leq t \cdot \hat{c}^T \boldsymbol{x}_{\boldsymbol{c}} + \phi(\boldsymbol{x}_{\boldsymbol{c}})$$

The different between $f_t(x_{\hat{c}})$ and $f_t(x_c)$ is

 $\begin{aligned} tc^T \boldsymbol{x}_{\hat{c}} + \boldsymbol{\phi}(\boldsymbol{x}_{\hat{c}}) - tc^T \boldsymbol{x}_c - \boldsymbol{\phi}(\boldsymbol{x}_c) \\ &\leq t(c^T \boldsymbol{x}_{\hat{c}} + \hat{c}^T \boldsymbol{x}_c - \hat{c}^T \boldsymbol{x}_{\hat{c}} - c^T \boldsymbol{x}_c) \\ &\leq 4tn2^{3L} \end{aligned}$

For $t = 1/2^{\Omega(L)}$) the last term becomes constant. Hence, using damped Newton we can move from $x_{\hat{c}}$ to x_c quickly.

In total for this analysis we require $\mathcal{O}(\sqrt{mL})$ outer iterations for the whole algorithm.

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