- inequalities  $Ax \leq b$ ;  $m \times n$  matrix A with rows  $a_i^T$
- $P = \{x \mid Ax \le b\}; P^{\circ} := \{x \mid Ax < b\}$
- interior point algorithm:  $x \in P^\circ$  throughout the algorithm
- for  $x \in P^\circ$  define

$$s_i(x) := b_i - a_i^T x$$

as the slack of the *i*-th constraint

logarithmic barrier function:

$$\phi(x) = -\sum_{i=1}^m \log(s_i(x))$$

Penalty for point *x*; points close to the boundary have a very large penalty.

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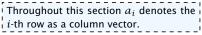
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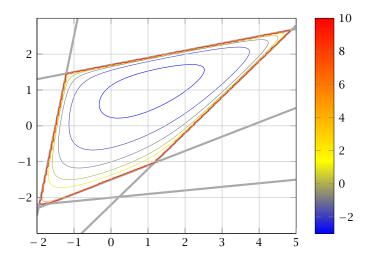
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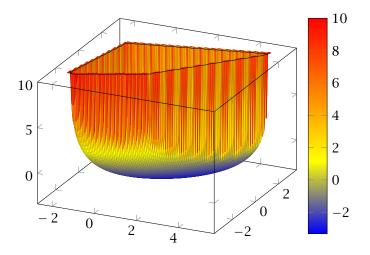
### **Penalty Function**





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### **Penalty Function**





### **Gradient and Hessian**

### **Taylor approximation:**

$$\phi(x+\epsilon) \approx \phi(x) + \nabla \phi(x)^T \epsilon + \frac{1}{2} \epsilon^T \nabla^2 \phi(x) \epsilon$$

Gradient:

$$\nabla \phi(x) = \sum_{i=1}^{m} \frac{1}{s_i(x)} \cdot a_i = A^T d_x$$

where  $d_x^T = (1/s_1(x), \dots, 1/s_m(x))$ . ( $d_x$  vector of inverse slacks)

Hessian:

$$H_{\mathbf{x}} := \nabla^2 \phi(\mathbf{x}) = \sum_{i=1}^m \frac{1}{s_i(\mathbf{x})^2} a_i a_i^T = A^T D_{\mathbf{x}}^2 A_i$$

with  $D_x = \operatorname{diag}(d_x)$ .

### **Gradient and Hessian**

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$$H_{x} := \nabla^{2} \phi(x) = \sum_{i=1}^{m} \frac{1}{s_{i}(x)^{2}} a_{i} a_{i}^{T} = A^{T} D_{x}^{2} A$$

with  $D_x = \text{diag}(d_x)$ .

### **Proof for Gradient**

$$\begin{split} \frac{\partial \phi(x)}{\partial x_i} &= \frac{\partial}{\partial x_i} \left( -\sum_r \ln(s_r(x)) \right) \\ &= -\sum_r \frac{\partial}{\partial x_i} \left( \ln(s_r(x)) \right) = -\sum_r \frac{1}{s_r(x)} \frac{\partial}{\partial x_i} \left( s_r(x) \right) \\ &= -\sum_r \frac{1}{s_r(x)} \frac{\partial}{\partial x_i} \left( b_r - a_r^T x \right) = \sum_r \frac{1}{s_r(x)} \frac{\partial}{\partial x_i} \left( a_r^T x \right) \\ &= \sum_r \frac{1}{s_r(x)} A_{ri} \end{split}$$

The *i*-th entry of the gradient vector is  $\sum_{r} 1/s_r(x) \cdot A_{ri}$ . This gives that the gradient is

$$\nabla \phi(x) = \sum_{r} 1/s_{r}(x)a_{r} = A^{T}d_{x}$$

### **Proof for Hessian**

$$\frac{\partial}{\partial x_j} \left( \sum_r \frac{1}{s_r(x)} A_{ri} \right) = \sum_r A_{ri} \left( -\frac{1}{s_r(x)^2} \right) \cdot \frac{\partial}{\partial x_j} \left( s_r(x) \right)$$
$$= \sum_r A_{ri} \frac{1}{s_r(x)^2} A_{rj}$$

Note that  $\sum_{r} A_{ri}A_{rj} = (A^{T}A)_{ij}$ . Adding the additional factors  $1/s_{r}(x)^{2}$  can be done with a diagonal matrix.

Hence the Hessian is

$$H_X = A^T D^2 A$$

 $H_X$  is positive semi-definite for  $x \in P^\circ$ 

 $u^{T}H_{x}u = u^{T}A^{T}D_{x}^{2}Au = ||D_{x}Au||_{2}^{2} \ge 0$ 

This gives that  $\phi(x)$  is convex.

If rank(A) = n,  $H_x$  is positive definite for  $x \in P^\circ$  $u^T H_x u = \|D_x A u\|_2^2 > 0$  for  $u \neq 0$ 

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 $E_{x} = \{ y \mid (y - x)^{T} H_{x} (y - x) \leq 1 \} = \{ y \mid ||y - x||_{H_{x}} \leq 1 \}$ 

Points in Ex are feasible!!!

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#### Points in *E<sub>x</sub>* are feasible!!!

$$(y - x)^{T} H_{x}(y - x) = (y - x)^{T} A^{T} D_{x}^{2} A(y - x)$$

$$= \sum_{i=1}^{m} \frac{(a_{i}^{T} (y - x))^{2}}{s_{i}(x)^{2}}$$

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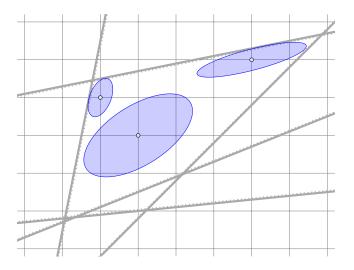
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# **Analytic Center**

 $x_{\mathrm{ac}} := \operatorname{arg\,min}_{x \in P^\circ} \phi(x)$ 

•  $x_{ac}$  is solution to

$$\nabla \phi(x) = \sum_{i=1}^{m} \frac{1}{s_i(x)} a_i = 0$$

- depends on the description of the polytope
- $x_{ac}$  exists and is unique iff  $P^{\circ}$  is nonempty and bounded



In the following we assume that the LP and its dual are strictly feasible and that rank(A) = n.

```
Central Path:
Set of points \{x^*(t) \mid t > 0\} with
```

 $x^*(t) = \operatorname{argmin}_x \{ t c^T x + \phi(x) \}$ 

- t = 0: analytic center
- $t = \infty$ : optimum solution

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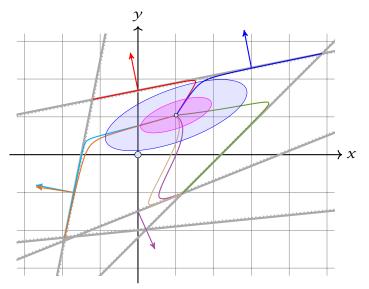
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# **Different Central Paths**





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#### Intuitive Idea:

Find point on central path for large value of t. Should be close to optimum solution.

### Questions:

- Is this really true? How large a t do we need?
- How do we find corresponding point  $x^*(t)$  on central path?



# The Dual

primal-dual pair:

#### Assumptions

primal and dual problems are strictly feasible;

•  $\operatorname{rank}(A) = n$ .

Note that the right LP in standard form is equal to  $\max\{-b^T y \mid -A^T y = c, x \ge 0\}$ . The dual of this is  $\min\{c^T x \mid -Ax \ge -b\}$  (variables x are unrestricted).

# **Force Field Interpretation**

Point  $x^*(t)$  on central path is solution to  $tc + \nabla \phi(x) = 0$ 

- We can view each constraint as generating a repelling force. The combination of these forces is represented by ∇φ(x).
- In addition there is a force tc pulling us towards the optimum solution.



### How large should *t* be?

Point  $x^*(t)$  on central path is solution to  $tc + \nabla \phi(x) = 0$ .

This means

$$tc + \sum_{i=1}^{m} \frac{1}{s_i(x^*(t))} a_i = 0$$

or

$$c + \sum_{i=1}^{m} z_i^*(t) a_i = 0$$
 with  $z_i^*(t) = \frac{1}{t s_i(x^*(t))}$ 

2000 is strictly dual feasible: (2022-5-0-0) 225-0)

if gap is less than  $1/2^{1010}$  we can snap to optimum point:

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*z*\*(*t*) is strictly dual feasible: (*A<sup>T</sup>z*\* + *c* = 0; *z*\* > 0)
 duality gap between *x* := *x*\*(*t*) and *z* := *z*\*(*t*) is

$$c^T x + b^T z = (b - Ax)^T z = \frac{m}{t}$$

If gap is less than 1/2<sup>Ω(L)</sup> we can snap to optimum point

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# How to find $x^*(t)$

#### First idea:

- start somewhere in the polytope
- use iterative method (Newtons method) to minimize  $f_t(x) := tc^T x + \phi(x)$



#### Quadratic approximation of $f_t$

$$f_t(x + \epsilon) \approx f_t(x) + \nabla f_t(x)^T \epsilon + \frac{1}{2} \epsilon^T H_{f_t}(x) \epsilon$$

Suppose this were exact:

$$f_t(x + \epsilon) = f_t(x) + \nabla f_t(x)^T \epsilon + \frac{1}{2} \epsilon^T H_{f_t}(x) \epsilon$$

Then gradient is given by:

$$\nabla f_t(x + \epsilon) = \nabla f_t(x) + H_{f_t}(x) \cdot \epsilon$$

Note that for the one-dimensional case  $g(\epsilon) = f(x) + f'(x)\epsilon + \frac{1}{2}f''(x)\epsilon^2$ , then  $g'(\epsilon) = f'(x) + f''(x)\epsilon$ .



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Observe that  $H_{f_t}(x) = H(x)$ , where H(x) is the Hessian for the function  $\phi(x)$  (adding a linear term like  $tc^T x$ does not affect the Hessian). Also  $\nabla f_t(x) = tc + \nabla \phi(x)$ .

We want to move to a point where this gradient is 0:

**Newton Step** at  $x \in P^{\circ}$ 

$$\Delta x_{\mathsf{nt}} = -H_{f_t}^{-1}(x)\nabla f_t(x) = -H_{f_t}^{-1}(x)(tc + \nabla \phi(x)) = -(A^T D_x^2 A)^{-1}(tc + A^T d_x)$$

**Newton Iteration:** 

 $x := x + \Delta x_{nt}$ 

## **Measuring Progress of Newton Step**

#### Newton decrement:

 $\lambda_t(x) = \|D_x A \Delta x_{\mathsf{nt}}\| \\ = \|\Delta x_{\mathsf{nt}}\|_{H_x}$ 

Square of Newton decrement is linear estimate of reduction if we do a Newton step:

 $-\lambda_t(x)^2 = \nabla f_t(x)^T \Delta x_{\mathsf{nt}}$ 

•  $\lambda_t(x) = 0$  iff  $x = x^*(t)$ 

•  $\lambda_t(x)$  is measure of proximity of x to  $x^*(t)$ 

Recall that  $\Delta x_{nt}$  fulfills  $-H(x)\Delta x_{nt} = \nabla f_t()$ .

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•  $\lambda_t(x)$  is measure of proximity of x to  $x^*(t)$ 

Recall that  $\Delta x_{nt}$  fulfills  $-H(x)\Delta x_{nt} = \nabla f_t()$ .

#### **Theorem 2** If $\lambda_t(x) < 1$ then

- $x_+ := x + \Delta x_{nt} \in P^\circ$  (new point feasible)
- $\lambda_t(x_+) \leq \lambda_t(x)^2$

This means we have quadratic convergence. Very fast.

#### feasibility:

►  $\lambda_t(x) = \|\Delta x_{nt}\|_{H_x} < 1$ ; hence  $x_+$  lies in the Dikin ellipsoid around x.

bound on  $\lambda_t(x^+)$ : we use  $D := D_x = \text{diag}(d_x)$  and  $D_+ := D_{x^+} = \text{diag}(d_{x^+})$ 

To see the last equality we use Pythagoras

 $||a||^2 + ||a + b||^2 = ||b||^2$ 

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$$a^{T}(a+b)$$

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If  $\lambda_t(x)$  is large we do not have a guarantee.

Try to avoid this case!!!



## **Path-following Methods**

Try to slowly travel along the central path.

Algorithm 1 PathFollowing

- 1: start at analytic center
- 2: while solution not good enough do
- 3: make step to improve objective function
- 4: recenter to return to central path

#### simplifying assumptions:

- a first central point  $x^*(t_0)$  is given
- $x^*(t)$  is computed exactly in each iteration

#### $\epsilon$ is approximation we are aiming for

start at  $t = t_0$ , repeat until  $m/t \le \epsilon$ 

- compute  $x^*(\mu t)$  using Newton starting from  $x^*(t)$
- ► *t* := µ*t*

where  $\mu = 1 + 1/(2\sqrt{m})$ 

gradient of  $f_{t^+}$  at ( $x = x^*(t)$ )

$$\nabla f_{t^+}(x) = \nabla f_t(x) + (\mu - 1)tc$$
$$= -(\mu - 1)A^T D_X \vec{1}$$

This holds because  $0 = \nabla f_t(x) = tc + A^T D_x \vec{1}$ .

The Newton decrement is

$$\begin{split} \lambda_{t^{+}}(x)^{2} &= \nabla f_{t^{+}}(x)^{T} H^{-1} \nabla f_{t^{+}}(x) \\ &= (\mu - 1)^{2} \vec{1}^{T} B (B^{T} B)^{-1} B^{T} \vec{1} \qquad B = D_{x}^{T} A \\ &\leq (\mu - 1)^{2} m \\ &= 1/4 \end{split}$$

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## **Number of Iterations**

the number of Newton iterations per outer tors. Since it is a projection maiteration is very small; in practise only 1 or  $2^{1} \frac{\text{trix}}{1000} (P^{2} = P)$  it can only have

#### Number of outer iterations:

We need  $t_k = \mu^k t_0 \ge m/\epsilon$ . This holds when

$$k \geq \frac{\log(m/(\epsilon t_0))}{\log(\mu)}$$

We get a bound of

$$\mathcal{O}\left(\sqrt{m}\log\frac{m}{\epsilon t_0}\right)$$

We show how to get a starting point with  $t_0 = 1/2^L$ . Together with  $\epsilon \approx 2^{-L}$  we get  $\mathcal{O}(L\sqrt{m})$  iterations.

EADS II Harald Räcke Explanation for previous slide  $P = B(B^TB)^{-1}B^T$  is a symmetric real-valued matrix; it has *n* linearly independent Eigenvectors. Since it is a projection matrix ( $P^2 = P$ ) it can only have Eigenvalues 0 and 1 (because the Eigenvalues of  $P^2$  are  $\lambda_i^2$ , where  $\lambda_i$  is Eigenvalue of *P*). The expression

$$\max_{v} \frac{v^T P v}{v^T v}$$

gives the largest Eigenvalue for P. Hence,  $\vec{1}^T P \vec{1} \le \vec{1}^T \vec{1} = m$ 

We assume that the polytope (not just the LP) is bounded. Then  $Av \leq 0$  is not possible.

For 
$$x \in P^\circ$$
 and direction  $v \neq 0$  define

$$\sigma_{\chi}(v) := \max_{i} \frac{a_{i}^{T} v}{s_{i}(\chi)}$$

 $a_i^T v$  is the change on the left hand side of the *i*-th constraint when moving in direction of v.

If  $\sigma_x(v) > 1$  then for one coordinate this change is larger than the slack in the constraint at position x.

By downscaling v we can ensure to stay in the polytope.

**Observation:** 

 $x + \alpha v \in P$  for  $\alpha \in \{0, 1/\sigma_x(v)\}$ 



Suppose that we move from x to  $x + \alpha v$ . The linear estimate says that  $f_t(x)$  should change by  $\nabla f_t(x)^T \alpha v$ .

The following argument shows that  $f_t$  is well behaved. For small  $\alpha$  the reduction of  $f_t(x)$  is close to linear estimate.

 $f_t(x + \alpha v) - f_t(x) = tc^T \alpha v + \phi(x + \alpha v) - \phi(x)$  $\phi(x + \alpha v) - \phi(x)$ 



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The following argument shows that  $f_t$  is well behaved. For small  $\alpha$  the reduction of  $f_t(x)$  is close to linear estimate.

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 $\phi(x + \alpha v) - \phi(x) = -\sum_{i} \log(s_i(x + \alpha v)) + \sum_{i} \log(s_i(x))$  $= -\sum_{i} \log(s_i(x + \alpha v)/s_i(x))$  $= -\sum_{i} \log(1 - a_i^T \alpha v/s_i(x))$ 



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Define  $w_i = a_i^T v / s_i(x)$  and  $\sigma = \max_i w_i$ . Then

 $f_t(x + \alpha v) - f_t(x) - \nabla f_t(x)^T \alpha v$ 

For 
$$|x| < 1$$
,  $x \le 0$ :  
 $x + \log(1 - x) = -\frac{x^2}{2} - \frac{x^3}{3} - \frac{x^4}{4} - \dots \ge -\frac{x^2}{2} = -\frac{y^2}{2} \frac{x^2}{y^2}$   
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## $\nabla f_t(x)^T \alpha v$ **Damped Newton Method** $= (tc^T + \sum_i a_i^T / s_i(x)) \alpha v$ $= tc^T \alpha v + \sum_i \alpha w_i$ Note that $||w|| = ||v||_{H_x}$ . Define $w_i = a_i^T v / s_i(x)$ and $\sigma = \max_i w_i$ . Then $f_t(x + \alpha v) - f_t(x) - \nabla f_t(x)^T \alpha v$ $= -\sum_{i} (\alpha w_i + \log(1 - \alpha w_i))$

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For 
$$x \ge 0$$
  
 $\frac{x^2}{2} \le \frac{x^2}{2} + \frac{x^3}{3} + \frac{x^4}{4} + \dots = -(x + \log(1 - x))$ 

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Damped Newton Iteration: In a damped Newton step we choose

$$x_{+} = x + \frac{1}{1 + \sigma_x(\Delta x_{\mathsf{nt}})} \Delta x_{\mathsf{nt}}$$

This means that in the above expressions we choose  $\alpha = \frac{1}{1+\alpha}$  and  $\nu = \Delta x_{nt}$ . Note that it wouldn't make sense to choose  $\alpha$  larger than 1 as this would mean that our real target  $(x + \Delta x_{nt})$  is inside the polytope but we overshoot and go further than this target.



#### 10 Karmarkars Algorithm

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#### Theorem:

In a damped Newton step the cost decreases by at least

 $\lambda_t(x) - \log(1 + \lambda_t(x))$ 

**Proof:** The decrease in cost is

$$-\alpha \nabla f_t(x)^T v + \frac{1}{\sigma^2} \|v\|_{H_x} (\alpha \sigma + \log(1 - \alpha \sigma))$$

Choosing  $\alpha = \frac{1}{1+\sigma}$  and  $v = \Delta x_{nt}$  gives

With  $v = \Delta x_{nt}$  we have  $||w||_2 = ||v||_{H_x} = \lambda_t(x)$ ; further recall that  $\sigma = ||w||_{\infty}$ ; hence  $\sigma \le \lambda_t(x)$ .

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The first inequality follows since the function  $\frac{1}{x^2}(x - \log(1+x))$  is monotonically decreasing.

 $\geq \lambda_t(x) - \log(1 + \lambda_t(x))$  $\geq 0.09$ 

#### for $\lambda_t(x) \ge 0.5$

**Centering Algorithm:** Input: precision  $\delta$ ; starting point *x* 

- **1.** compute  $\Delta x_{nt}$  and  $\lambda_t(x)$
- **2.** if  $\lambda_t(x) \leq \delta$  return x
- **3.** set  $x := x + \alpha \Delta x_{nt}$  with

$$\alpha = \begin{cases} \frac{1}{1 + \sigma_x(\Delta x_{\text{nt}})} & \lambda_t \ge 1/2 \\ 1 & \text{otw.} \end{cases}$$



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## Centering

#### Lemma 3

# The centering algorithm starting at $x_0$ reaches a point with $\lambda_t(x) \le \delta$ after

$$\frac{f_t(x_0) - \min_{\mathcal{Y}} f_t(\mathcal{Y})}{0.09} + \mathcal{O}(\log \log(1/\delta))$$

iterations.

This can be very, very slow...



# Let $P = \{Ax \le b\}$ be our (feasible) polyhedron, and $x_0$ a feasible point.

We change  $b \to b + \frac{1}{\lambda} \cdot \vec{1}$ , where  $L = \langle A \rangle + \langle b \rangle + \langle c \rangle$  (encoding length) and  $\lambda = 2^{2L}$ . Recall that a basis is feasible in the old LP iff it is feasible in the new LP.



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#### **Lemma** [without proof] The inverse of a matrix M can be represented with rational numbers that have denominators $z_{ij} = det(M)$ .

For two basis solutions  $x_B$ ,  $x_{\bar{B}}$ , the cost-difference  $c^T x_B - c^T x_{\bar{B}}$ can be represented by a rational number that has denominator  $z = \det(A_B) \cdot \det(A_{\bar{B}}) \cdot \lambda$ .

This means that in the perturbed LP it is sufficient to decrease the duality gap to  $1/2^{4L}$  (i.e.,  $t \approx 2^{4L}$ ). This means the previous analysis essentially also works for the perturbed LP.

For a point x from the polytope (not necessarily BFS) the objective value  $\bar{c}^T x$  is at most  $n2^M 2^L$ , where  $M \leq L$  is the encoding length of the largest entry in  $\bar{c}$ .



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Start at  $x_0$ .Note that an entry in  $\hat{c}$  fulfills  $|\hat{c}_i| \le 2^{2L}$ .This holds since the slack in every constraint<br/>at  $x_0$  is at least  $\lambda = 1/2^{2L}$ , and the gradient<br/>is the vector of inverse slacks.

 $x_0 = x^*(1)$  is point on central path for  $\hat{c}$  and t = 1.

You can travel the central path in both directions. Go towards 0 until  $t \approx 1/2^{\Omega(L)}$ . This requires  $O(\sqrt{m}L)$  outer iterations.

Let  $x_{\hat{c}}$  denote this point.

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Clearly,

$$t \cdot \hat{c}^T \boldsymbol{x}_{\hat{c}} + \phi(\boldsymbol{x}_{\hat{c}}) \leq t \cdot \hat{c}^T \boldsymbol{x}_{\boldsymbol{c}} + \phi(\boldsymbol{x}_{\boldsymbol{c}})$$

#### The different between $f_t(x_{\hat{c}})$ and $f_t(x_c)$ is

 $\begin{aligned} tc^T \boldsymbol{x}_{\hat{c}} + \boldsymbol{\phi}(\boldsymbol{x}_{\hat{c}}) - tc^T \boldsymbol{x}_c - \boldsymbol{\phi}(\boldsymbol{x}_c) \\ &\leq t(c^T \boldsymbol{x}_{\hat{c}} + \hat{c}^T \boldsymbol{x}_c - \hat{c}^T \boldsymbol{x}_{\hat{c}} - c^T \boldsymbol{x}_c) \\ &\leq 4tn2^{3L} \end{aligned}$ 

For  $t = 1/2^{\Omega(L)}$ ) the last term becomes constant. Hence, using damped Newton we can move from  $x_{\hat{c}}$  to  $x_c$  quickly.

In total for this analysis we require  $\mathcal{O}(\sqrt{mL})$  outer iterations for the whole algorithm.

Clearly,

$$t \cdot \hat{c}^T \boldsymbol{x}_{\hat{c}} + \phi(\boldsymbol{x}_{\hat{c}}) \leq t \cdot \hat{c}^T \boldsymbol{x}_{\boldsymbol{c}} + \phi(\boldsymbol{x}_{\boldsymbol{c}})$$

The different between  $f_t(x_{\hat{c}})$  and  $f_t(x_c)$  is

 $\begin{aligned} tc^T x_{\hat{c}} + \phi(x_{\hat{c}}) - tc^T x_c - \phi(x_c) \\ &\leq t(c^T x_{\hat{c}} + \hat{c}^T x_c - \hat{c}^T x_{\hat{c}} - c^T x_c) \\ &\leq 4tn 2^{3L} \end{aligned}$ 

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The different between  $f_t(x_{\hat{c}})$  and  $f_t(x_c)$  is

$$tc^{T}x_{\hat{c}} + \phi(x_{\hat{c}}) - tc^{T}x_{c} - \phi(x_{c})$$

$$\leq t(c^{T}x_{\hat{c}} + \hat{c}^{T}x_{c} - \hat{c}^{T}x_{\hat{c}} - c^{T}x_{c})$$

$$\leq 4tn2^{3t}$$

For  $t = 1/2^{\Omega(L)}$ ) the last term becomes constant. Hence, using damped Newton we can move from  $x_{\hat{c}}$  to  $x_c$  quickly.

In total for this analysis we require  $\mathcal{O}(\sqrt{m}L)$  outer iterations for the whole algorithm.

Clearly,

$$t \cdot \hat{c}^T \boldsymbol{x}_{\hat{c}} + \boldsymbol{\phi}(\boldsymbol{x}_{\hat{c}}) \leq t \cdot \hat{c}^T \boldsymbol{x}_c + \boldsymbol{\phi}(\boldsymbol{x}_c)$$

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$$\leq t(c^{T}x_{\hat{c}} + \hat{c}^{T}x_{c} - \hat{c}^{T}x_{\hat{c}} - c^{T}x_{c})$$

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$$\leq t(c^{T}x_{\hat{c}} + \hat{c}^{T}x_{c} - \hat{c}^{T}x_{\hat{c}} - c^{T}x_{c})$$

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