Duality

How do we get an upper bound to a maximization LP?

Note that a lower bound is easy to derive. Every choice of $a, b \ge 0$ gives us a lower bound (e.g. a = 12, b = 28 gives us a lower bound of 800).

If you take a conic combination of the rows (multiply the *i*-th row with $y_i \ge 0$) such that $\sum_i y_i a_{ij} \ge c_j$ then $\sum_i y_i b_i$ will be an upper bound.

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Duality

Lemma 3

The dual of the dual problem is the primal problem.

Proof:

- $w = \min\{b^T y \mid A^T y \ge c, y \ge 0\}$
- $w = -\max\{-b^T y \mid -A^T y \le -c, y \ge 0\}$

The dual problem is

- $z = -\min\{-c^T x \mid -Ax \ge -b, x \ge 0\}$
- $z = \max\{c^T x \mid Ax \le b, x \ge 0\}$

Duality

Definition 2

Let $z = \max\{c^T x \mid Ax \le b, x \ge 0\}$ be a linear program P (called the primal linear program).

The linear program D defined by

$$w = \min\{b^T y \mid A^T y \ge c, y \ge 0\}$$

is called the dual problem.

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Weak Duality Let $z = \max\{c^T x \mid Ax \le b, x \ge 0\}$ and $w = \min\{b^T y \mid A^T y \ge c, y \ge 0\}$ be a primal dual pair. x is primal feasible iff $x \in \{x \mid Ax \le b, x \ge 0\}$ y is dual feasible, iff $y \in \{y \mid A^T y \ge c, y \ge 0\}$. Theorem 4 (Weak Duality) Let \hat{x} be primal feasible and let \hat{y} be dual feasible. Then $c^T \hat{x} \le z \le w \le b^T \hat{y}$.

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Weak Duality

 $A^T \hat{y} \ge c \Rightarrow \hat{x}^T A^T \hat{y} \ge \hat{x}^T c \ (\hat{x} \ge 0)$

 $A\hat{x} \le b \Rightarrow y^T A \hat{x} \le \hat{y}^T b \ (\hat{y} \ge 0)$

This gives

$$c^T \hat{x} \le \hat{y}^T A \hat{x} \le b^T \hat{y}$$

Since, there exists primal feasible \hat{x} with $c^T \hat{x} = z$, and dual feasible \hat{y} with $b^T y = w$ we get $z \le w$.

If P is unbounded then D is infeasible.

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5.1 Weak Duality

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