Shortest Path

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\text{s.t.} & \forall S \in \mathcal{S} & \sum_{e\in\delta(S)} x_e & \geq & 1 \\
& \forall e \in E & x_e & \in & \{0,1\}
\end{array}
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$\mathcal{S}$ is the set of subsets that separate s from t.

The Dual:

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The Separation Problem for the Shortest Path LP is the Minimum Cut Problem.

18 Cuts & Metrics

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Observations:

Suppose that ℓ_e-values are solution to Minimum Cut LP.

- We can view ℓ_e as defining the length of an edge.
- Define $d(u,v) = \min_{\text{path } P \text{ btw. } u \text{ and } v} \sum_{e \in P} \ell_e$ as the Shortest Path Metric induced by ℓ_e.
- We have $d(u,v) = \ell_e$ for every edge $e = (u,v)$, as otw. we could reduce ℓ_e without affecting the distance between s and t.

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How do we round the LP?

- Let $B(s,r)$ be the ball of radius r around s (w.r.t. metric d). Formally:

$$B = \{v \in V \mid d(s,v) \le r\}$$

- For $0 \le r < 1$, $B(s,r)$ is an s-t-cut.

Which value of r should we choose? choose randomly!!!

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choose r u.a.r. (uniformly at random) from interval $[0,1)$

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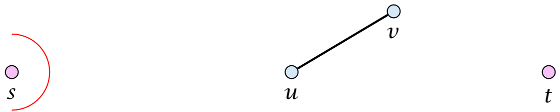

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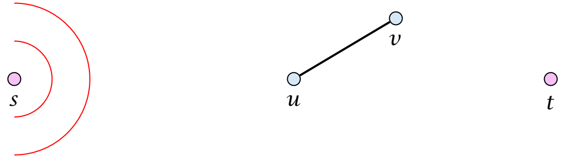

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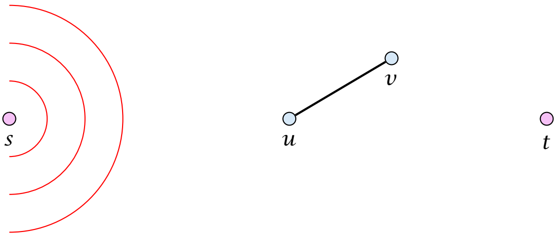

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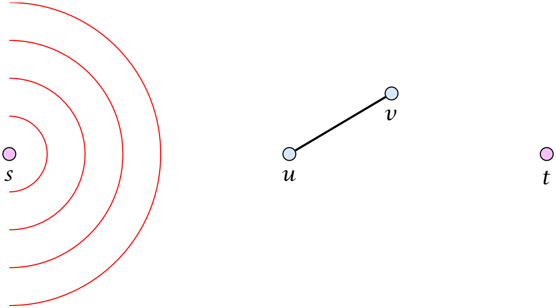

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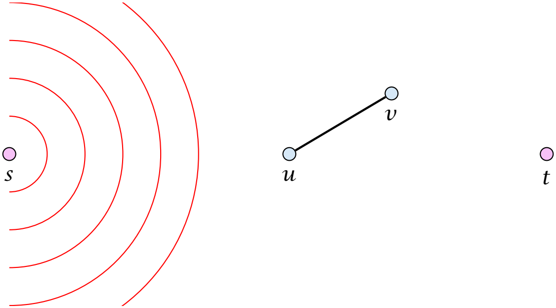

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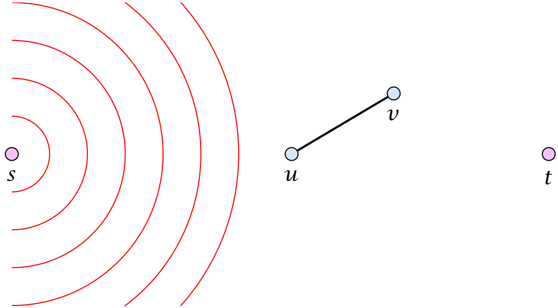

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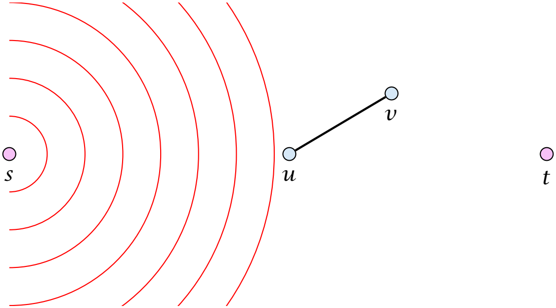

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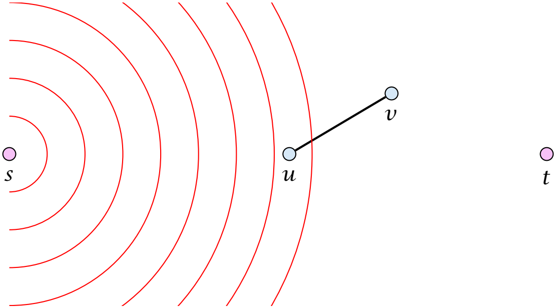

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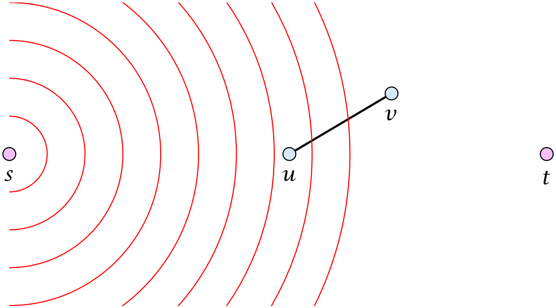

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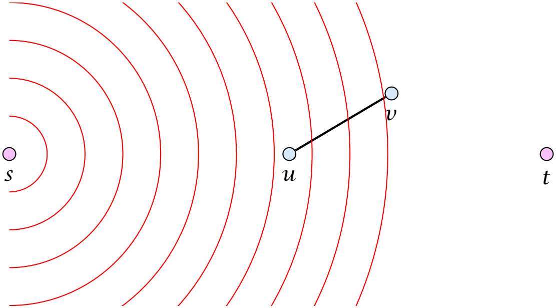

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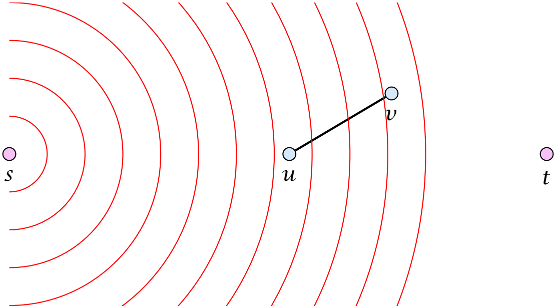

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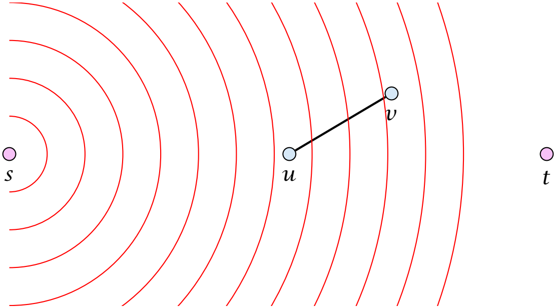

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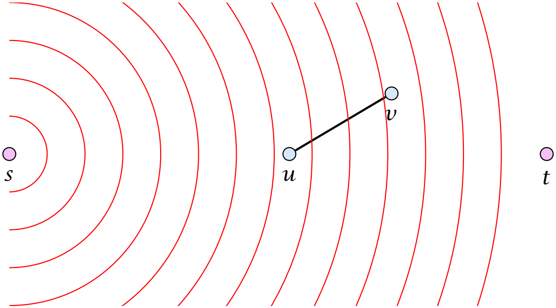

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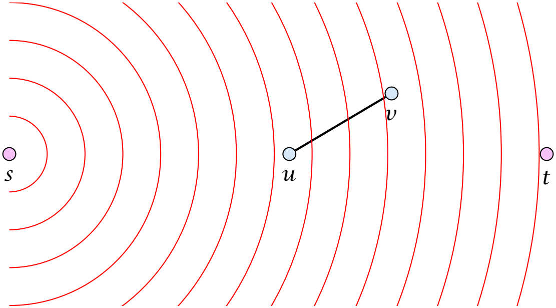

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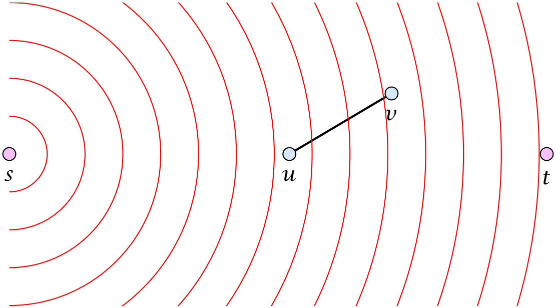

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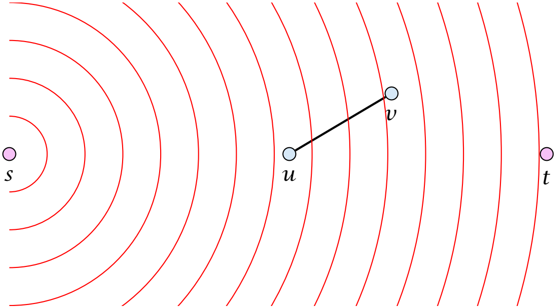

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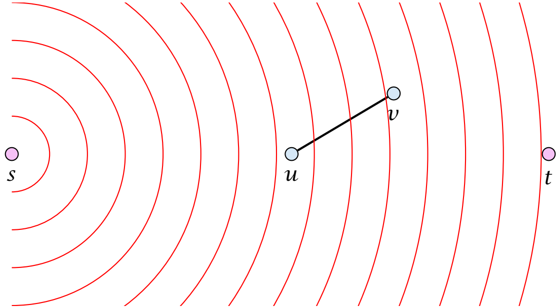

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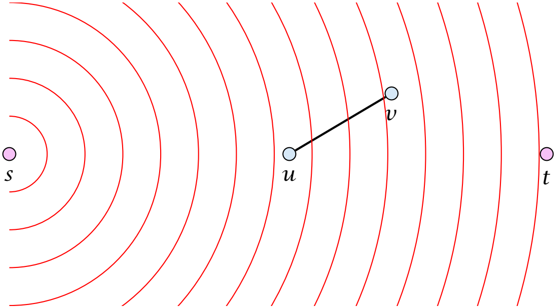

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$$\leq \ell_e$$

What is the expected size of a cut?

$$\mathrm{E}[\text{size of cut}] = \mathrm{E}[\textstyle\sum_e c(e)\,\Pr[e \text{ is cut}]]$$
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On the other hand:

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as the ℓ_e are the solution to the Mincut LP *relaxation*.

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Given a graph $G = (V, E)$, together with source-target pairs s_i, t_i, $i = 1, \dots, k$, and a capacity function $c : E \to \mathbb{R}^+$ on the edges. Find a subset $F \subseteq E$ of the edges such that all s_i-t_i pairs lie in different components in $G = (V, E \setminus F)$.

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Re-using the analysis for the single-commodity case is difficult.

$$\Pr[e \text{ is cut}] \le \textbf{?}$$

- If for some R the balls $B(s_i, R)$ are disjoint between different sources, we get a $1/R$ approximation.
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- Replace the graph G by a graph G', where an edge of length ℓ_e is replaced by ℓ_e/δ edges of length δ.
- Let $B(s_i, z)$ be the ball in G' that contains nodes v with distance $d(s_i, v) \le z\delta$.

Algorithm 1 RegionGrowing(s_i, p)

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z ← 0
repeat
    flip a coin (Pr[heads] = p)
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while $\exists s_i$-t_i pair in G' **do**
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- if it reaches the edge then it either cuts the edge or protects the edge from being cut by other sources
- if we choose $p = \delta$ the probability of cutting an edge is only its LP-value; our expected cost are at most OPT.

Problem:

We may not cut all source-target pairs.

A component that we remove may contain an s_i-t_i pair.

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- choose $p = 6\ln k \cdot \delta$
- we make $\frac{1}{2\delta}$ trials before reaching radius $1/2$.
- we say a Region Growing is not successful if it does not terminate before reaching radius $1/2$.

$$\Pr[\text{not successful}] \leq (1-p)^{\frac{1}{2\delta}} = \left((1-p)^{1/p}\right)^{\frac{p}{2\delta}} \leq e^{-\frac{p}{2\delta}} \leq \frac{1}{k^3}$$

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What is expected cost?

$$\begin{aligned}\mathrm{E}[\text{cutsize}] = {} & \Pr[\text{success}] \cdot \mathrm{E}[\text{cutsize} \mid \text{success}] \\ & + \Pr[\text{no success}] \cdot \mathrm{E}[\text{cutsize} \mid \text{no success}]\end{aligned}$$

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Note: success means all source-target pairs separated

We assume $k \ge 2$.

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We assume $k \ge 2$.

If we are not successful we simply perform a trivial k-approximation.

This only increases the expected cost by at most $\frac{1}{k^2} \cdot k\mathrm{OPT} \le \mathrm{OPT}/k$.

Hence, our final cost is $\mathcal{O}(\ln k) \cdot \mathrm{OPT}$ in expectation.