Duality

How do we get an upper bound to a maximization LP?

Note that a lower bound is easy to derive. Every choice of $a, b \ge 0$ gives us a lower bound (e.g. a = 12, b = 28 gives us a lower bound of 800).

If you take a conic combination of the rows (multiply the i-th row with $y_i \ge 0$) such that $\sum_i y_i a_{ij} \ge c_i$ then $\sum_i y_i b_i$ will be an upper bound.



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Duality

Lemma 31

The dual of the dual problem is the primal problem.

Proof:

$$w = -\max\{-b^T y \mid -A^T y \le -c, y \ge 0\}$$

The dual problem is

$$z = -\min\{-c^T x \mid -Ax \ge -b, x \ge 0\}$$

$$z = \max\{c^T x \mid Ax \le b, x \ge 0\}$$

Duality

Definition 30

Let $z = \max\{c^T x \mid Ax \le b, x \ge 0\}$ be a linear program P (called the primal linear program).

The linear program D defined by

$$w = \min\{b^T \gamma \mid A^T \gamma \ge c, \gamma \ge 0\}$$

is called the dual problem.

Harald Räcke

5.1 Weak Duality

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Weak Duality

Let $z = \max\{c^T x \mid Ax \le b, x \ge 0\}$ and $w = \min\{b^T \gamma \mid A^T \gamma \ge c, \gamma \ge 0\}$ be a primal dual pair.

x is primal feasible iff $x \in \{x \mid Ax \le b, x \ge 0\}$

 γ is dual feasible, iff $\gamma \in \{\gamma \mid A^T \gamma \geq c, \gamma \geq 0\}$.

Theorem 32 (Weak Duality)

Let \hat{x} be primal feasible and let \hat{y} be dual feasible. Then

$$c^T \hat{x} \leq z \leq w \leq b^T \hat{v}$$
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$$A^T \hat{y} \ge c \Rightarrow \hat{x}^T A^T \hat{y} \ge \hat{x}^T c \ (\hat{x} \ge 0)$$

$$A\hat{x} \le b \Rightarrow y^T A \hat{x} \le \hat{y}^T b \ (\hat{y} \ge 0)$$

This gives

$$c^T \hat{x} \leq \hat{y}^T A \hat{x} \leq b^T \hat{y}$$
.

Since, there exists primal feasible \hat{x} with $c^T\hat{x}=z$, and dual feasible \hat{y} with $b^T\hat{y}=w$ we get $z\leq w$.

If P is unbounded then D is infeasible.



5.1 Weak Duality

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