
Efficient Algorithms and Data Structures I

Deadline: None, Tutorial exercises only.

Tutorial Exercise 1

For constants $c > 0$, $0 < \varepsilon < 1/2$ and $k > 1$, arrange the following functions of n in non-decreasing asymptotic order so that $f_i(n) \in O(f_{i+1}(n))$ for two consecutive functions in your sequence. Also indicate whether $f_i(n) \in \Theta(f_{i+1}(n))$ holds or not.

$$n^k, n^{1+\sin(n)}, \log(n!), n^{k+\varepsilon}, n^n, n, n^k(\log n)^c, n!, 2^n, 3^n, n \log \log n, n \log(n), n^\varepsilon, n^{1/\log n}$$

Use any base for the logarithm you like the most. The natural logarithm facilitates the calculations.

Tutorial Exercise 2

Show by using the basic definition of the Θ -notation, that for any real constants a and b , where $b > 0$,

$$(n + a)^b = \Theta(n^b)$$

Tutorial Exercise 3

Jonathan is frustrated. His boss has given him a list of integers a_1, \dots, a_n with a weird assignment: For each $i \in [n]$, compute the product $b_i = \prod_{j \in [n], j \neq i} a_j$. Unfortunately, the latest update of his operating system has broken the division operator on his machine.

- At first, Jonathan's logarithm operator is still working well (in $\mathcal{O}(1)$ time). Find a way for him to compute all b_i 's in $\mathcal{O}(n)$ steps.
- The next update to his system breaks the logarithm operator as well. Help Jonathan find a new $\mathcal{O}(n)$ time algorithm to compute all b_i 's.

The advanced reader who skips parts that appear too elementary may miss more than the less advanced reader who skips parts that appear too complex.

- G. Pòlya